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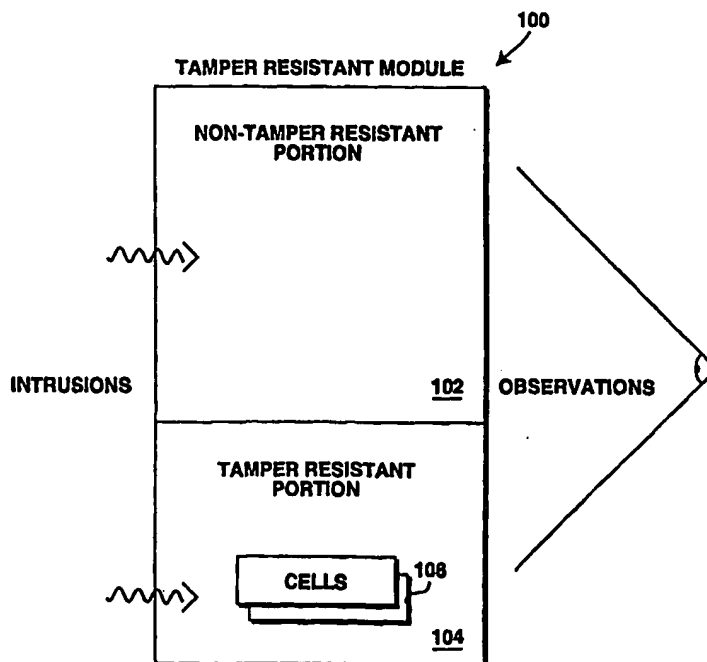
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(54) Title: A TAMPER RESISTANT PLAYER FOR SCRAMBLED CONTENTS

(57) Abstract

In one apparatus, a group of plain text and obfuscated cells (106) of programming instructions is provided to implement a descrambler that descrambles scrambled content to generate descrambled content. In another apparatus, a group of plain text and obfuscated cells of programming instructions is provided to implement an authenticator that provides appropriate authentication challenges to a scrambled content provider, and generates appropriate authentication responses to authentication challenges from the scrambled content provider. In yet another apparatus, a group of plain text and obfuscated cells of programming instructions is provided to implement an integrity verifier that performs integrity verification on a decoder. In yet another apparatus, a group of plain text and obfuscated cells of programming instructions is provided to implement a secrets holder that holds a number of secrets associated with playing scrambled contents.



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A Tamper Resistant Player For Scrambled Contents

RELATED APPLICATIONS

This application is a continuation-in-part application to U.S. Patent Application, number 08/662,679, filed on June 13, 1996, entitled Tamper Resistant Methods and Apparatus, and to U.S. Patent Application, number <to be assigned>, filed on August 6, 1997, entitled Cell Array Providing Non-Persistent Secret Storage Through A Mutation Cycle (Express Mail No. EM531554811US).

BACKGROUND OF THE INVENTION

1. Field of the Invention

The present invention relates to the field of system security. More specifically, the present invention relates to a tamper resistant player for scrambled contents.

2. Background Information

Content management, such as management of scrambled DVD contents, require the basic integrity of the management operations to be assumed, or at least verified. While a number of security approaches such as encryption and decryption techniques are known in the art, unfortunately, the security approaches can be readily compromised, because these applications and the security approaches are implemented on

systems with an open and accessible architecture, that renders both hardware and software including the security approaches observable and modifiable by a malevolent user or a malicious program.

Thus, a system based on open and accessible architecture is a fundamentally insecure platform, notwithstanding the employment of security measures. However, openness and accessibility offer a number of advantages, contributing to these systems' successes. Therefore, what is required are techniques that will render the operations of a scrambled content player, such as a DVD player, virtually unobservable or unmodifiable on these fundamentally insecure platforms, notwithstanding their openness and accessibility.

SUMMARY OF THE INVENTION

In one apparatus, a group of plain text and obfuscated cells of programming instructions is provided to implement a descrambler that descrambles scrambled content to generate descrambled content.

In another apparatus, a group of plain text and obfuscated cells of programming instructions is provided to implement an authenticator that provides appropriate authentication challenges to a scrambled content provider, and generates appropriate authentication responses to authentication challenges from the scrambled content provider.

In yet another apparatus, a group of plain text and obfuscated cells of programming instructions is provided to implement an integrity verifier that performs integrity verification on a decoder.

In yet another apparatus, a group of plain text and obfuscated cells of programming instructions is provided to implement a secrets holder that holds a number of secrets associated with playing scrambled contents. In one embodiment, the secrets include secrets used in a mutual authentication process, and the secrets used for descrambling the scrambled content.

BRIEF DESCRIPTION OF DRAWINGS

The present invention will be described by way of embodiments, but not limitations, illustrated in the accompanying drawings in which like references denote similar elements, and in which:

Figure 1 is a block diagram illustrating an overview of an exemplary tamper resistant module incorporated with various teachings of the present invention;

Figures 2 - 3 are two flow charts illustrating one embodiment each of the operational flows, at start-up time and during runtime, of an integrity verification method of the present invention

Figure 4 is a flow chart illustrating one embodiment of the operational flow of an intruder detection method of the present invention;

Figures 5 - 6 are two flow charts illustrating one embodiment each of the operational flows of two observation detection methods of the present invention;

Figure 7 is a block diagram illustrating one embodiment of a coupling technique of the present invention for inter-coupling various tamper resistant methods;

Figure 8 is a block diagram illustrating one embodiment of a tamper resistant player for scrambled contents, incorporated with the teachings of the present invention; and

Figure 9 is a block diagram illustrating one embodiment of a computer system suitable for practicing the present invention.

DETAILED DESCRIPTION OF THE INVENTION

In the following description, various aspects of the present invention will be described. However, it will be apparent to those skilled in the art that the present invention may be practiced with only some or all aspects of the present invention. For purposes of explanation, specific numbers, materials and configurations are set forth in order to provide a thorough understanding of the present invention. However, it will also be apparent to one skilled in the art that the present invention may be practiced without the specific details. In other instances, well known features are omitted or simplified in order not to obscure the present invention.

Parts of the description will be presented in terms of operations performed by a computer system, using terms such as data, flags, bits, values, characters, strings, numbers and the like, consistent with the manner commonly employed by those skilled in the art to convey the substance of their work to others skilled in the art. As well understood by those skilled in the art, these quantities take the form of electrical, magnetic, or optical signals capable of being stored, transferred, combined, and otherwise manipulated through mechanical and electrical components of the computer system; and the term computer system include general purpose as well as special purpose data processing machines, systems, and the like, that are standalone, adjunct or embedded.

Various operations will be described as multiple discrete steps in turn in a manner that is most helpful in understanding the present invention, however, the order of description should not be construed as to imply that these operations are necessarily order dependent, in particular, the order of presentation.

Referring now to **Figure 1**, wherein a block diagram illustrating one embodiment of an exemplary tamper resistant module incorporated with the various teachings of the present invention is shown. As illustrated, exemplary tamper resistant module **100** includes non-tamper resistant portion **102**, and tamper resistant portion **104**. For the illustrated embodiment, the two portions are linked together to form a single executable module. For the purpose of this application, the term module is used in a general

sense to mean a structural relationship between the various portions that facilitates exclusive communications between the portions.

As described in the parent application, number 08/662,679, non-tamper resistant portion **102** includes a number of plain text programming instructions implementing various non-sensitive services of exemplary tamper resistant module **100**, whereas tamper resistant portion **104** includes various groups of plain text and obfuscated cells **106** of programming instructions implementing various sensitive services of exemplary tamper resistant module **100**. Each group of cells that implements a sensitive service or a collection of sensitive services includes at least one plain text cell **106**. Briefly, the secrets associated with the services are distributed in time and space, and obfuscated. The number of obfuscated cells employed to obfuscate a service is service or sensitivity dependent. Generally, the larger number of obfuscated cells employed, the more difficult it will be for the obfuscation to be "decoded". For a more detailed description, see parent application, number 08/662,679.

Additionally, in accordance with the present invention, selected groups of plain text and obfuscated cells **106** incorporate a number of tamper resistant measures to verify during operation that exemplary tamper resistant module **100** has not been intruded nor being observed. The number of groups employing these tamper resistant measures, as well as the frequencies and the number of tamper resistant measures employed are also service or sensitivity dependent. As will be described in more details below, these tamper

resistant measures include a number of integrity verification measures and a number of anti-observation measures. The integrity verification measures include first integrity verification measure that verifies the integrity of non-tamper resistant portion **102** during run time as well as start-up time, and a second integrity verification measure that verifies an invocation of a group of plain text and obfuscated cells is not originated from an intruder. The anti-observation measures include a first anti-observation measure that verifies the processor executing module **100** is not operating in a mode that supports single step execution, and a second anti-observation measure that verifies elapsed execution times are consistent with normal unobserved execution.

Figures 2 - 3 illustrate one embodiment of the operational flow of the first integrity verification measure. **Figure 2** illustrates the operational flow at start-up time, whereas **Figure 3** illustrates the operational flow during run time. As shown in **Fig. 2**, at start-up time, for the illustrated embodiment, a group of cells (GOC) incorporated with this first integrity verification measure scans non-tamper resistant portion **102** and calculates a signature for non-tamper resistant portion **102**, block **108**. Next, for the illustrated embodiment, the GOC retrieves a signature pre-stored for non-tamper resistant portion **102**, block **110**. The GOC then compares the two signatures to verify the generated signature, blocks **112 - 114**. If the generated signature is successfully verified, meaning that non-tamper resistant portion **102** has not been modified, the GOC allows the start-up process to continue, without skipping any verification dependent operations, block **116**, otherwise, the GOC

causes the start-up process to continue, skipping the verification dependent operations, block 118. An example of verification dependent operations is operations associated with setting up the secrets required for delivering certain sensitive services.

As shown in **Figure 3**, at a verification check time during run time, for the illustrated embodiment, a GOC incorporated with this first integrity verification measure scans a next portion of non-tamper resistant portion 102 and incrementally calculates a signature for non-tamper resistant portion 102, block 120. The GOC then updates the signature being incrementally calculated, block 122. Next, the GOC checks if the end of non-tamper resistant portion 102 has been reached, block 124. If the end has not been reached, the process terminates, otherwise the process continues at block 126.

At block 126, the GOC retrieves a signature pre-stored for non-tamper resistant portion 102, block 126. The GOC then compares the two signatures to verify the generated signature, blocks 128 – 130. If the generated signature is successfully verified, meaning that non-tamper resistant portion 102 has not been modified, the GOC allows execution of module 100 to continue, otherwise, the GOC causes execution of module 100 to terminate, block 132. Causing module to terminate may be achieved in any number of ways known in the art. Depending on the application, it may be preferable to cause the module to fail further downstream from the point the non-tamper resistant portion's integrity failed verification.

In other words, the run time integrity check is performed incrementally over a number of verification check times during an execution run. Those skilled in the art will appreciate the incremental approach is particularly useful for performance sensitive services. The number of verification check times employed for an execution run is service or sensitivity dependent.

Figure 4 illustrates one embodiment of the operational flow of the second integrity verification measure. At invocation time, for the illustrated embodiment, a GOC incorporated with this second integrity verification measure retrieves a return address for the invocation, block **134**. For the illustrated embodiment, the GOC determines if the return address is within the address space of module **100**, block **136**. If the return address is within the address space of module **100**, meaning that the invocation did not originate from an intruder, the GOC allows execution of module **100** to continue, block **138**, otherwise, the GOC causes execution of module **100** to terminate, block **140**. Similarly, causing module to terminate may be achieved in any number of ways known in the art. Depending on the application, it may be preferable to cause the module to fail further downstream from the point the intrusion is detected.

Figure 5 illustrates one embodiment of the operational flow of the first anti-observation measure. At a pre-selected point in time during an execution run, for the illustrated embodiment, a GOC incorporated with this first anti-observation measure retrieves a

processor execution mode state variable, block 142. For the illustrated embodiment, the GOC determines if the state variable denotes an execution mode that supports single step execution, e.g. a debug mode, block 144. If the state variable denotes an execution mode that does not support single step execution, meaning that execution of module 100 is not being observed, the GOC allows execution of module 100 to continue, block 146, otherwise, the GOC causes execution of module 100 to terminate, block 148. Similarly, causing module to terminate may be achieved in any number of ways known in the art. Depending on the application, it may be preferable to cause the module to fail further downstream from the point observation is detected. The number of times as well as the precise points in time during an execution run where the processor's execution mode is checked is service or sensitivity dependent.

Figure 6 illustrates one embodiment of the operational flow of the second anti-observation measure. At a pre-selected point in time during an execution run, for the illustrated embodiment, a GOC incorporated with this second anti-observation measure retrieves a timer value from the processor executing module 100, and record the retrieved timer value (timestamp), block 150. The GOC then continues to perform the normal services it is designed to provide, block 152. At a pre-selected later point in time, the GOC checks an amount of elapsed execution time since the last timestamp to determine if the amount of elapsed execution has exceeded a predetermined threshold, blocks 154 – 156. If the elapsed execution time does not exceed the predetermined

threshold, meaning that execution of module 100 is not being observed (e.g. by setting breakpoints), the GOC allows execution of module 100 to continue, block 158, otherwise, the GOC causes execution of module 100 to terminate, block 160. Similarly, causing module to terminate may be achieved in any number of ways known in the art. Depending on the application, it may be preferable to cause the module to fail further downstream from the point observation is detected. The number of times as well as the precise points in time during an execution run where the amount of elapsed execution time since a last timestamp is checked is service or sensitivity dependent.

Figure 7 illustrates one embodiment of a coupling technique for inter-coupling tamper resistant measures. As illustrated, the different tamper resistant measures are inter-coupled by having the measures share a common storage location, e.g. in memory, for key values associated with the various tamper resistant measures. For the illustrated embodiment, a GOC stores a key for retrieving secrets in portion 162 of storage location 168, and a timestamp for determining whether execution of module 100 is being observed in storage location 168 less portion 162. In determining elapsed execution time, the GOC only employs the bits higher than portion 162. Additionally, the GOC uses lower order bits 164 as a seed to generate the pseudo random numbers employed in an authentication process. Thus, if an intruder attempts to modify the timestamp to defeat the elapsed execution time check measure, it will cause the authentication process as well as any attempt to retrieve secrets to fail. Similarly, if an intruder attempts to modify the

seed for generating pseudo random number to defeat the authentication process, it will cause the elapsed execution time check as well as any attempt to retrieve secrets to fail.

Figure 8 illustrates one embodiment of a tamper resistant player for scrambled content applying the tamper resistant teachings of the present invention. As shown, for the illustrated embodiment, tamper resistant player **170** includes non-tamper resistant components **171** and tamper resistant decoder **172**. Non-tamper resistant components **171** are intended to represent a broad category of general service components, such as end user interfaces. These general service components may provide any one of a number of variety of services, implemented using any one of a number of variety of techniques known in the art. Tamper resistant decoder **172** receives scrambled compressed content, and in response, descrambles as well as decompresses the content to output appropriate signals to render the content, e.g. YUV video and AC3 audio.

Tamper resistant decoder **172** includes non-tamper resistant portion **175**, tamper resistant portion **174**, **176**, **178** and **180**, and signature **173** for non-tamper resistant portion **175**. Non-tamper resistant portion **175** is constituted with plain text programming instructions, whereas tamper resistant portion **174**, **176**, **178** and **180** is constituted with multiple groups of plain text and obfuscated cells of programming instructions. Non-tamper resistant portion **175** and tamper resistant portion **174**, **176**, **178** and **180**, including signature **173**, are structurally related to facilitate

exclusive communication between the portions. For the illustrated embodiment, the two portions are linked together as a single executable module.

Non-tamper resistant portion 175 selectively invokes the services of integrated tamper resistant portion 174, 176, 178 and 180 to effectuate descrambling of the scrambled content, including causing player 170 and a scrambled content provider device to be mutually authenticated with one another. Non-tamper resistant portion 175 decompresses the unscrambled compressed content to generate the above described output signals. Signature 173 is pre-stored in a predetermined location to facilitate start-up time and run time integrity verification as described earlier.

For the illustrated embodiment, tamper resistant services of tamper resistant decoder 172 includes tamper resistant descrambler 174 for receiving scrambled content, and in response, descrambling the scrambled content to generate the descrambled content for non-tamper resistant portion of decoder 172. In one embodiment, tamper resistant descrambler 174 employs secret keys retrieved from tamper resistant secrets holder 180 to descramble the scrambled content. The number of secret keys employed, and the nature of the keys are application dependent, and they are not essential to the understanding of the present invention. Tamper resistant descrambler 174 is constituted with a group of plain text and obfuscated cells of programming instructions. In one embodiment, the core descrambling service is disposed in a plain text cell to provide enhanced performance. In one embodiment, the

GOC is equipped with the above described intruder detection integrity verification measure and the single step execution mode detection anti-observation measure. In one embodiment, the GOC is also equipped with the elapsed execution time detection anti-observation measure. In one embodiment, the GOC is equipped with multiple ones of the elapsed execution time detection anti-observation measure. In one embodiment, the elapsed execution time detection anti-observation measure is also inter-coupled with the process for retrieving the secret keys associated with descrambling scrambled content, and the authentication process for mutually authenticating player 170 and a scrambled content provider device.

For the illustrated embodiment, tamper resistant services of tamper resistant decoder 172 also includes tamper resistant authenticator 176 for authenticating tamper resistant player 170 to a scrambled content provider device and to authenticate the scrambled content provider device to tamper resistant player 170. In one embodiment, tamper resistant authenticator 176 employs secret keys retrieved from tamper resistant secrets holder 180 to conduct the authentication process. The number of secret keys employed, and the nature of the keys are application dependent, and they are not essential to the understanding of the present invention. In one embodiment, tamper resistant authenticator 176 is constituted with a group of plain text and obfuscated cells of programming instructions. In one embodiment, the GOC is equipped with the above described intruder detection integrity verification measure, and the single step execution mode detection anti-observation measure. In one

embodiment, the GOC is also equipped with the elapsed execution time detection anti-observation measure. In one embodiment, the GOC is equipped with multiple ones of the elapsed execution time detection anti-observation measures. In one embodiment, the elapsed execution time detection anti-observation measure is also inter-coupled with the process for retrieving the secret keys associated with descrambling scrambled content, and the authentication process for mutually authenticating player 170 and a scrambled content provider device.

For the illustrated embodiment, tamper resistant services of tamper resistant decoder 172 also includes tamper resistant integrity verifier 178 for integrity verifying non-tamper resistant portion of decoder 172 at start-up time, and during run time. In one embodiment, tamper resistant integrity verifier 178 provides secret keys to be employed for mutually authenticating player 170 and a scrambled content provider device to secrets holder 180. The number of secret keys employed, and the nature of the keys are application dependent, and they are not essential to the understanding of the present invention. In one embodiment, tamper resistant integrity verifier 178 is constituted with a group of plain text and obfuscated cells of programming instructions. In one embodiment, the GOC is equipped with the single step execution mode detection anti-observation measure. In one embodiment, the GOC is also equipped with the elapsed execution time detection anti-observation measure. In one embodiment, the GOC is equipped with multiple ones of the elapsed execution time detection anti-observation measures. In one embodiment, the elapsed execution

time detection anti-observation measure is also inter-coupled with the authentication process for retrieving the secret keys associated with descrambling scrambled content, and the authentication process for mutually authenticating player 170 and a scrambled content provider device.

Lastly, as alluded to, for the illustrated embodiment, tamper resistant services of tamper resistant decoder 172 includes tamper resistant secrets holder 180 for storing secrets associated with descrambling scrambled content. Secrets holder 180 also stores secrets associated with an authentication process for authenticating tamper resistant player 170 to a scrambled content provider device and to authenticate the scrambled content provider device to tamper resistant player 170. In one embodiment, tamper resistant secrets holder 180 is constituted with a group of plain text and obfuscated cells of programming instructions in a cell array form as described in parent application, number <to be assigned> (Express mail number EM531554811US). In one embodiment, the GOC is equipped with the above described intruder detection integrity verification measure, and the single step execution mode detection anti-observation measure. In one embodiment, the GOC is also equipped with the elapsed execution time detection anti-observation measure. In one embodiment, the GOC is equipped with multiple ones of the elapsed execution time detection anti-observation measures.

Thus, even if player 170 receives its content inputs through an "open" bus, the content is nevertheless protected, as the

content will be provided to player 170 over the "open" bus in scrambled form. Furthermore, the secrets associated with descrambling the scrambled content, as well as the programming instructions performing the descrambling are protected from intrusion as well as from observation. Yet, performance sensitive operations, such as the core descrambling service, are not burdened. Lastly, the tamper resistant services, i.e. descrambler 174, authenticator 176 etc. are highly portable, and may be linked up with any number of decoder implementations.

Figure 9 illustrates one embodiment of a computer system suitable for practicing the present invention. As shown, for the illustrated embodiment, computer system 200 includes processor 202, processor bus 206, high performance I/O bus 210 and standard I/O bus 220. Processor bus 206 and high performance I/O bus 210 are bridged by host bridge 208, whereas I/O buses 210 and 212 are bridged by I/O bus bridge 212. Coupled to processor bus 206 is cache 204. Coupled to high performance I/O bus 210 are system memory 214 and video memory 216, to which video display 218 is coupled. Coupled to standard I/O bus 220 are disk drive 222, keyboard and pointing device 224 and DVD-ROM 226.

These elements perform their conventional functions known in the art. In particular, disk drive 222 and system memory 214 are used to store a permanent and a working copy of the tamper resistant application of the present invention, when executed by processor 202. The permanent copy may be pre-loaded into disk

drive **222** in factory, loaded from a distribution medium (not shown), or down loaded from on-line/networked distribution source (not shown). The constitutions of these elements are known. Any one of a number of implementations of these elements known in the art may be used to form computer system **200**.

Of course, computer systems of alternate constitutions, including computer systems of alternate architectures may also be employed to practice the present invention.

In general, while the present invention have been described in terms of the above illustrated embodiments, those skilled in the art will recognize that the invention is not limited to the embodiments described. The present invention can be practiced with modification and alteration within the spirit and scope of the appended claims. The description is thus to be regarded as illustrative instead of restrictive on the present invention.

Thus, a tamper resistant player for scrambled contents has been described.

CLAIMS

What is claimed is:

1. An apparatus comprising:
a storage medium having stored therein a first group of plain text and obfuscated cells of programming instructions designed to implement a descrambler that descrambles scrambled content to generate descrambled content; and
an execution unit coupled to the storage medium for executing the programming instructions.
2. The apparatus as set forth in claim 1, wherein the first group of plain text and obfuscated cells of programming instructions include an obfuscated cell of programming instructions designed to implement a frequently invoked descrambling subfunction of the descrambler.
3. The apparatus as set forth in claim 1, wherein the first group of plain text and obfuscated cells of programming instructions is equipped to verify that an invocation of the descrambler is not originated from an intruder, and/or execution of the first group of plain text and obfuscated cells of programming instructions is not being observed.
4. The apparatus as set forth in claim 1, wherein the storage medium further having stored therein a second group of plain text and obfuscated cells of programming instructions designed to

implement an authenticator that provides appropriate authentication challenges to a scrambled content provider device, and generates appropriate authentication responses to authentication challenges from the scrambled content provider device.

5. The apparatus as set forth in claim 4, wherein the second group of plain text and obfuscated cells of programming instructions is equipped to verify that execution of the second group of plain text and obfuscated cells of programming instructions is not being observed.

6. The apparatus as set forth in claim 1, wherein the storage medium further having stored therein a second group of plain text and obfuscated cells of programming instructions designed to implement a secrets holder that holds a plurality of secrets associated with playing the scrambled contents.

7. The apparatus as set forth in claim 6, wherein the second group of plain text and obfuscated cells of programming instructions is equipped to verify that execution of the second group of plain text and obfuscated cells of programming instructions is not being observed.

8. The apparatus as set forth in claim 1, wherein the storage medium further having stored therein a plurality of plain text programming instructions designed to decode the descrambled content, the plurality of plain text programming instructions having a

structural relationship with the descrambler that facilitates exclusive communication of the descrambled content from the descrambler.

9. The apparatus as set forth in claim 8, wherein the storage medium further having stored therein a second group of plain text and obfuscated cells of programming instructions designed to implement an integrity verifier that performs integrity verification on the plurality of plain text programming instructions.

10. The apparatus as set forth in claim 9, wherein the integrity verifier is designed to perform the integrity verification on the plurality of plain text programming instructions at start-up time, and/or incrementally at run time.

11. The apparatus as set forth in claim 9, wherein the second group of plain text and obfuscated cells of programming instructions is equipped to verify that execution of the second group of plain text and obfuscated cells of programming instructions is not being observed.

12. An apparatus comprising:
a storage medium having stored therein a plurality of plain text programming instructions designed to decode descrambled content, and a first group of plain text and obfuscated cells of programming instructions designed to implement an integrity verifier that performs integrity verification on the plurality of plain text programming instructions; and

an execution unit coupled to the storage medium for executing the programming instructions.

13. The apparatus as set forth in claim 12, wherein the integrity verifier is designed to perform the integrity verification on the plain text programming instructions at start-up time, and/or incrementally at run time.

14. The apparatus as set forth in claim 12, wherein the storage medium further having stored therein a second group of plain text and obfuscated cells of programming instructions designed to implement an authenticator that provides appropriate authentication challenges to a scrambled content provider device, and generates appropriate authentication responses to authentication challenges from the scrambled content provider device.

15. The apparatus as set forth in claim 14, wherein the second group of plain text and obfuscated cells of programming instructions is equipped to verify that execution of the second group of plain text and obfuscated cells of programming instructions is not being observed.

16. The apparatus as set forth in claim 12, wherein the storage medium further having stored therein a second group of plain text and obfuscated cells of programming instructions designed to implement a secrets holder that holds a plurality of secrets associated with playing the scrambled contents.

17. The apparatus as set forth in claim 16, wherein the second group of plain text and obfuscated cells of programming instructions is equipped to verify that execution of the second group of plain text and obfuscated cells of programming instructions is not being observed.

18. An apparatus comprising:

a storage medium having stored therein a first group of plain text and obfuscated cells of programming instructions designed to implement a secrets holder that holds a plurality of secrets associated with playing scrambled contents; and

an execution unit coupled to the storage medium for executing the programming instructions.

19. The apparatus as set forth in claim 18, wherein the first group of plain text and obfuscated cells of programming instructions is equipped to verify that execution of the first group of plain text and obfuscated cells of programming instructions is not being observed.

20. The apparatus as set forth in claim 19, wherein the storage medium further having stored therein a second group of plain text and obfuscated cells of programming instructions that implement an authenticator that provides appropriate authentication challenges to a scrambled content provider device, and generates appropriate authentication responses to authentication challenges from the scrambled content provider device.

21. The apparatus as set forth in claim 20, wherein the second group of plain text and obfuscated cells of programming instructions is equipped to verify that execution of the second group of plain text and obfuscated cells of programming instructions is not being observed.

22. An apparatus comprising:

a storage medium having stored therein a first group of plain text and obfuscated cells of programming instructions designed to implement an authenticator that provides appropriate authentication challenges to a scrambled content provider device, and generates appropriate authentication responses to authentication challenges from the scrambled content provider device; and

an execution unit coupled to the storage medium for executing the programming instructions.

23. The apparatus as set forth in claim 22, wherein the second group of plain text and obfuscated cells of programming instructions is equipped to verify that execution of the second group of plain text and obfuscated cells of programming instructions is not being observed.

24. A method comprising:

a) executing a first group of plain text and obfuscated cells of programming instructions to perform integrity verification on a plurality of plain text programming instructions designed to decode descrambled content; and

b) executing a second group of plain text and obfuscated cells of programming instructions to store a plurality of secrets associated with descrambling scrambled content if integrity of the plurality of plain text programming instructions is verified in (a).

25. The method of claim 24, wherein each of (a) and (b) includes determining if execution of the first /second group of plain text and obfuscated cells of programming instructions is being observed.

26. The method of claim 24, wherein the method further comprises (c) executing a third group of plain text and obfuscated cells of programming instructions to provide appropriate authentication challenges to a scrambled content provider device, and to provide appropriate responses to authentication challenges of the scrambled content provider device.

27. The method of claim 26, wherein each of (a), (b) and (c) includes determining if execution of the first /second/third group of plain text and obfuscated cells of programming instructions is being observed.

28. The method of claim 24, wherein the method further comprises (d) executing a fourth group of plain text and obfuscated cells of programming instructions to descramble scrambled content to generate the descrambled content if authentication with the scrambled content provider device is successful in (c).

29. The method of claim 28, wherein each of (a), (b), (c) and (d) includes determining if execution of the first /second/third/fourth group of plain text and obfuscated cells of programming instructions is being observed.

30. The method of claim 28, wherein (d) includes determining if an invocation of the fourth group of programming instructions is originated from an intruder.

31. The method of claim 28, wherein the method further includes (e) executing the plurality of plain text programming instructions to decode the descrambled content.

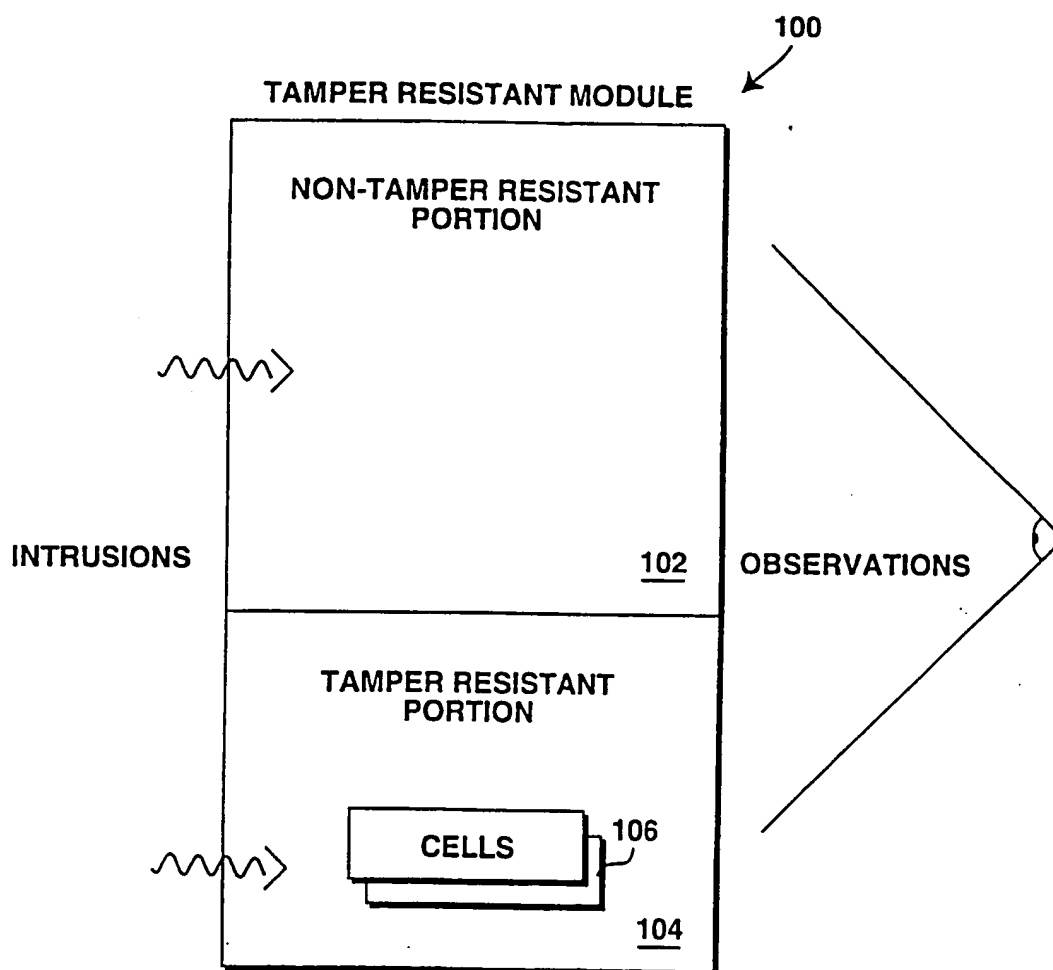


Fig. 1

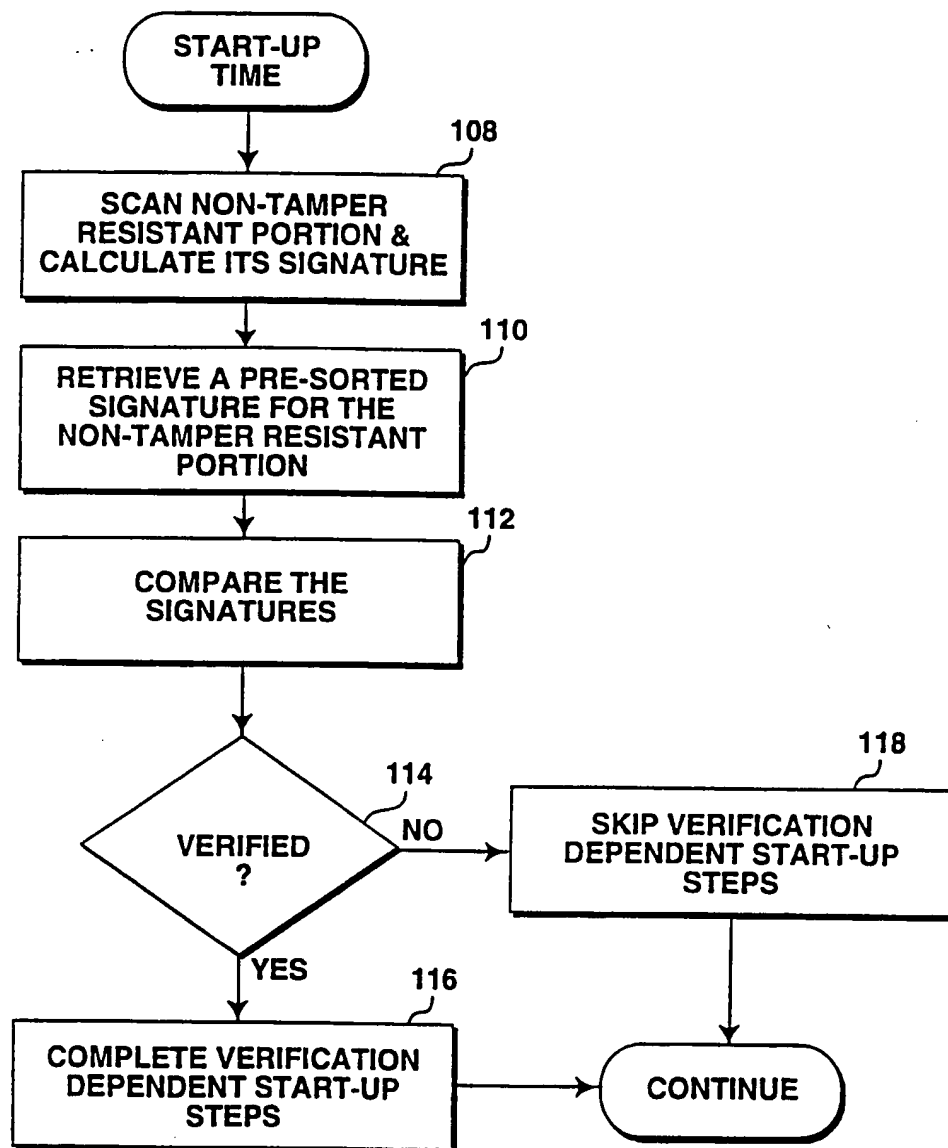
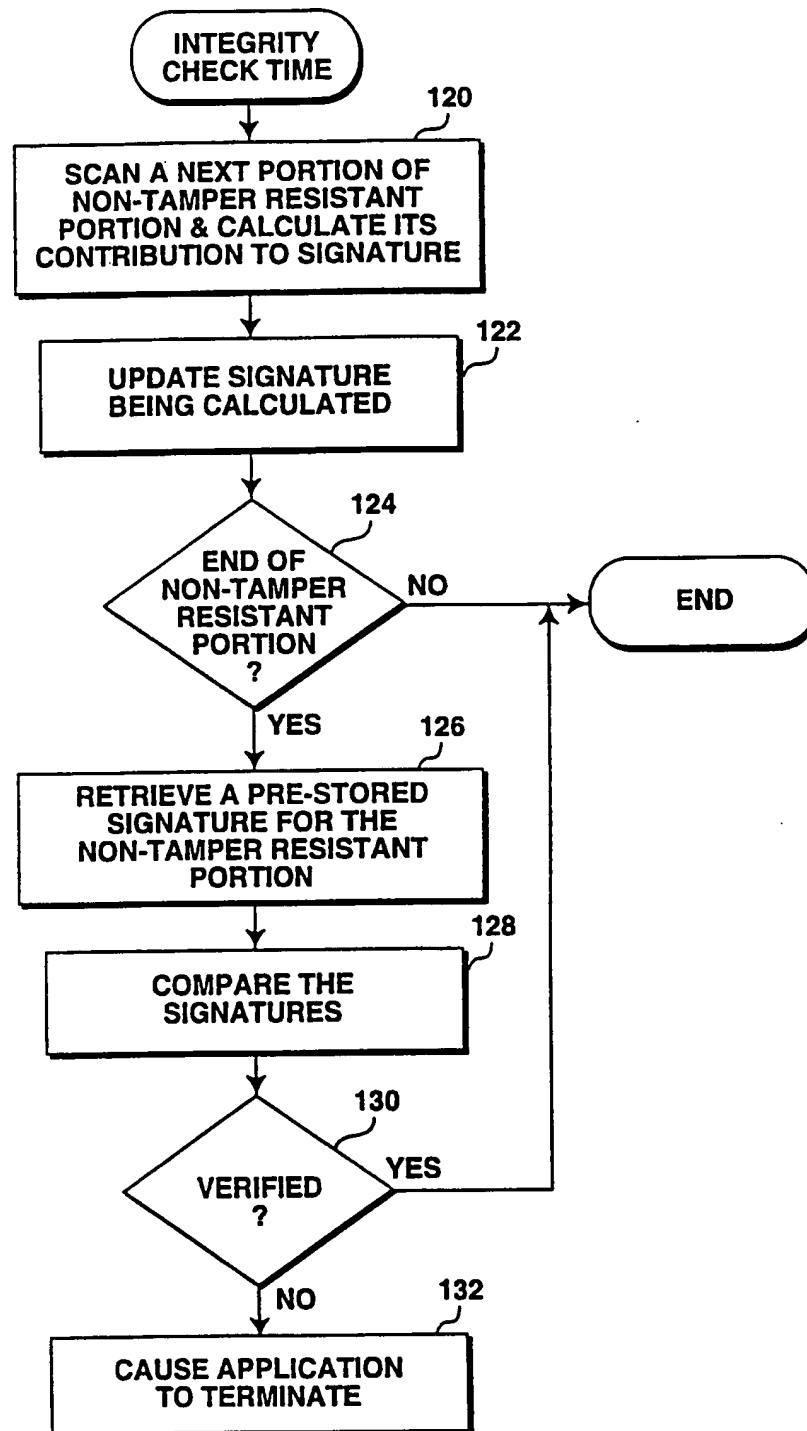
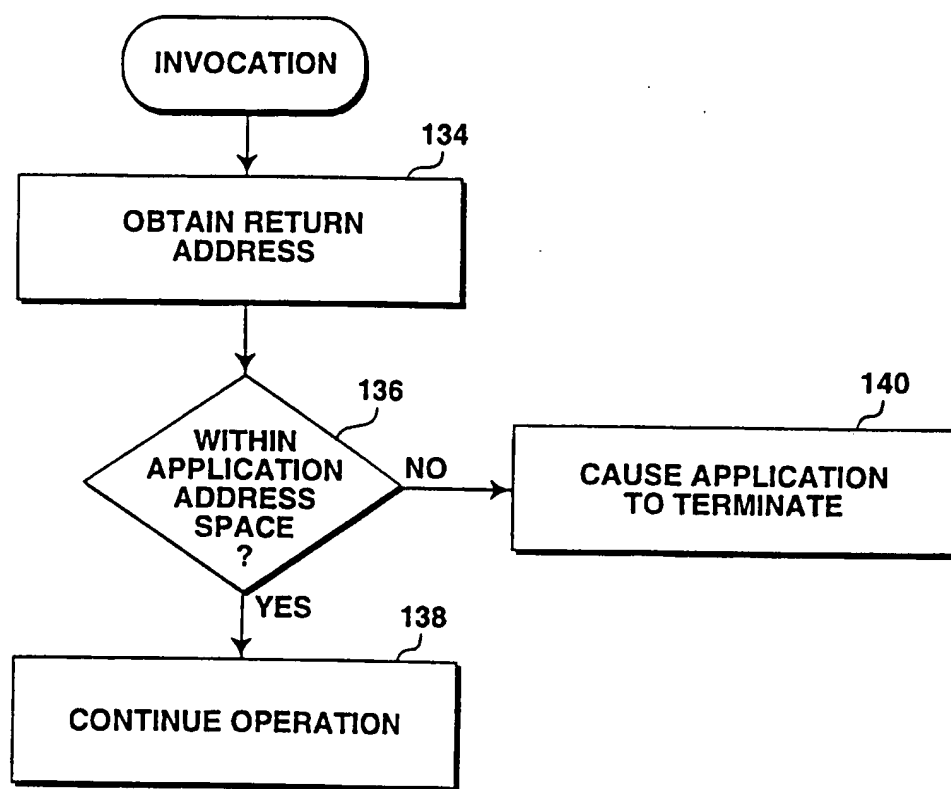
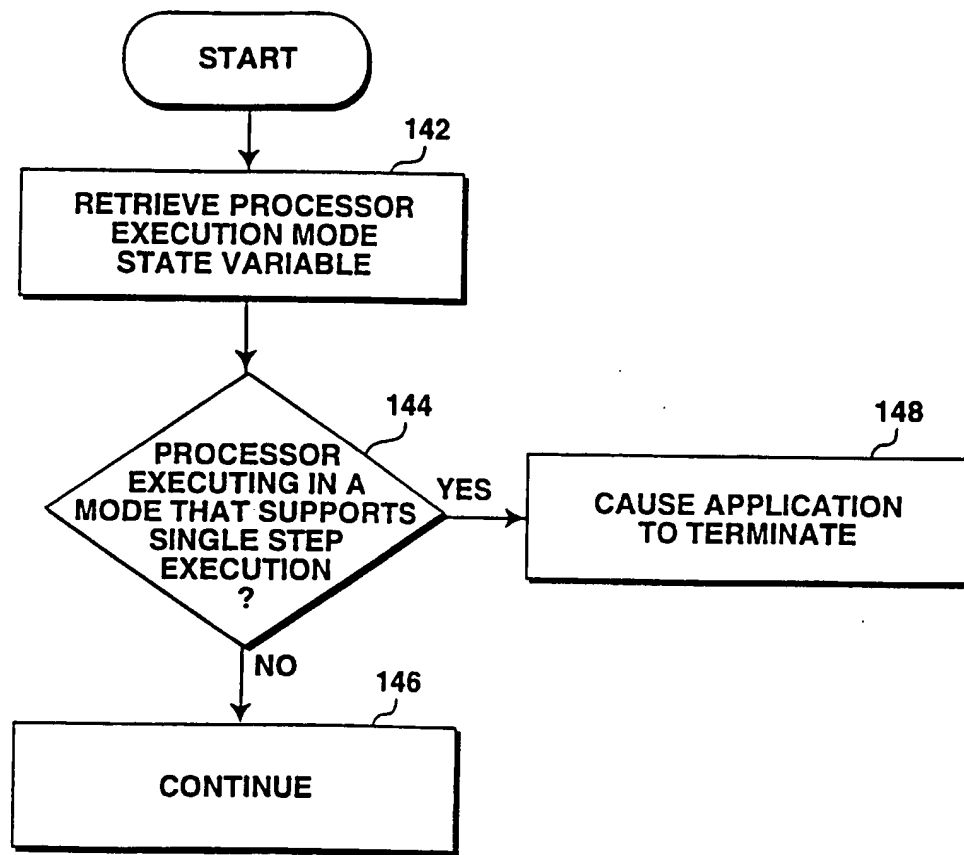
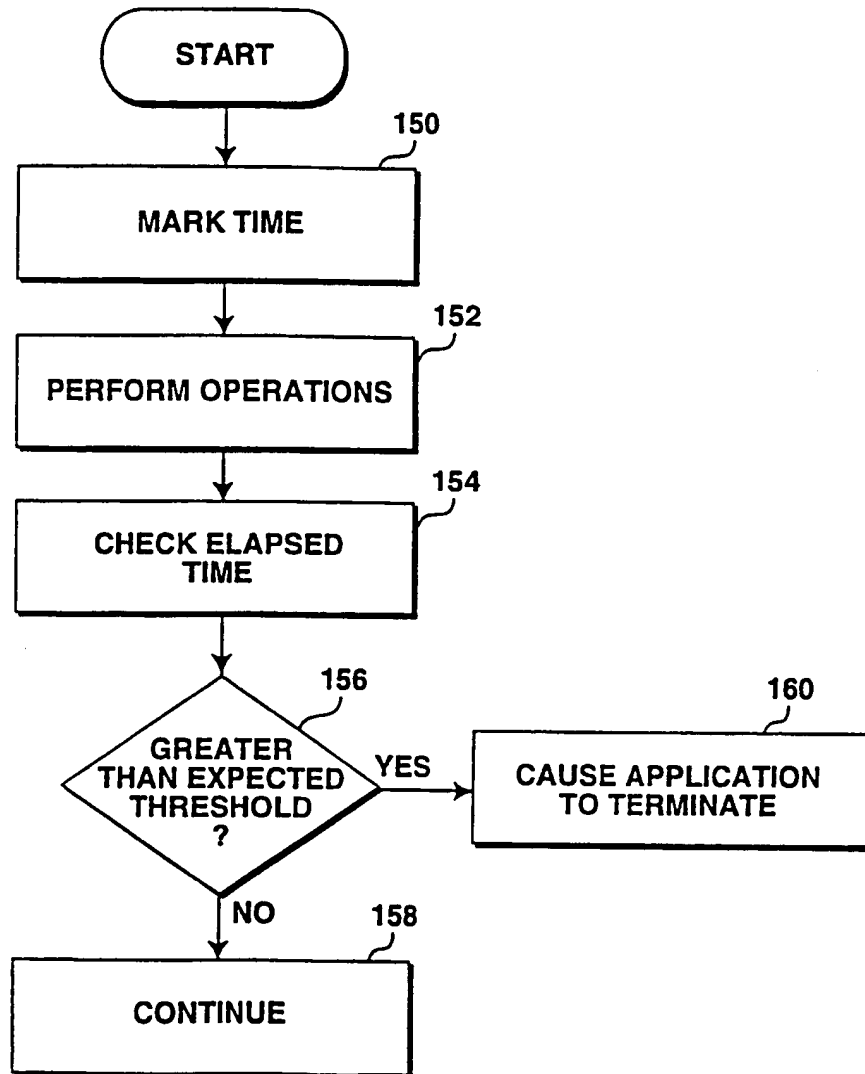


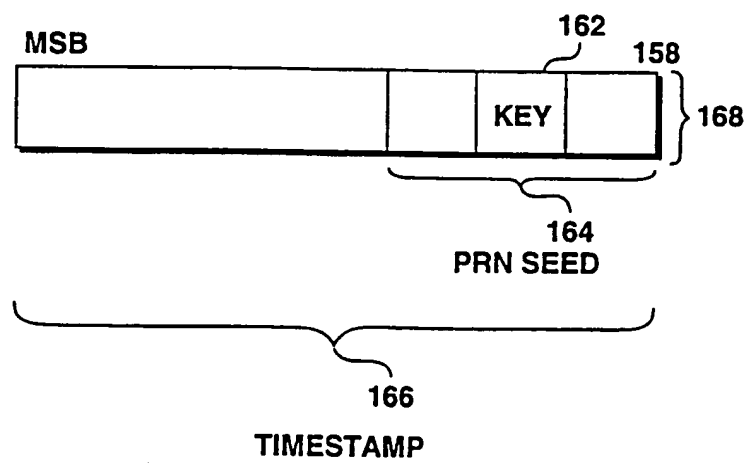
Fig. 2

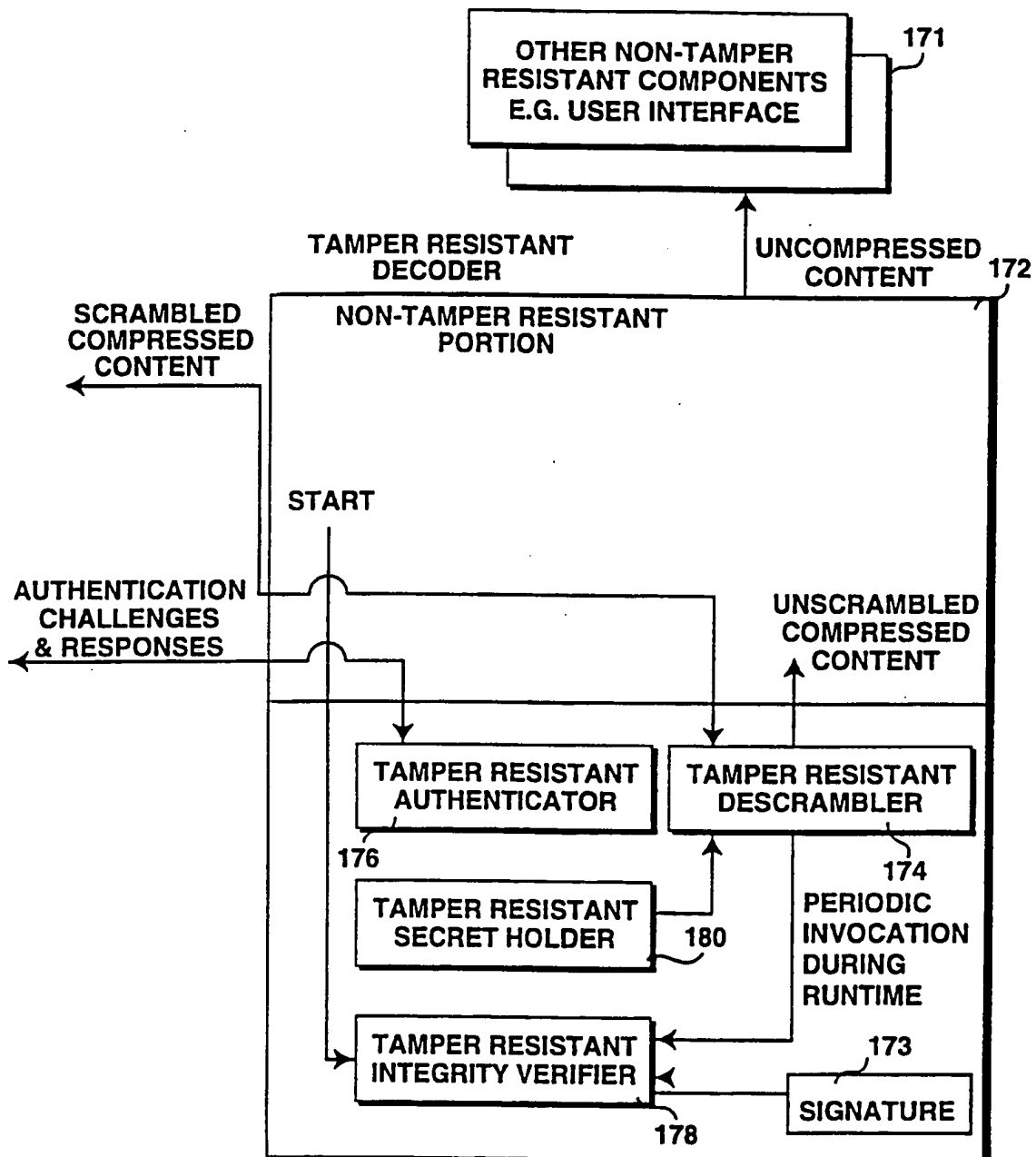
**Fig. 3**

**Fig. 4**

**Fig. 5**

**Fig. 6**

**Fig. 7**



TAMPER RESISTANT PLAYER FOR SCRAMBLED CONTENT

170

Fig. 8

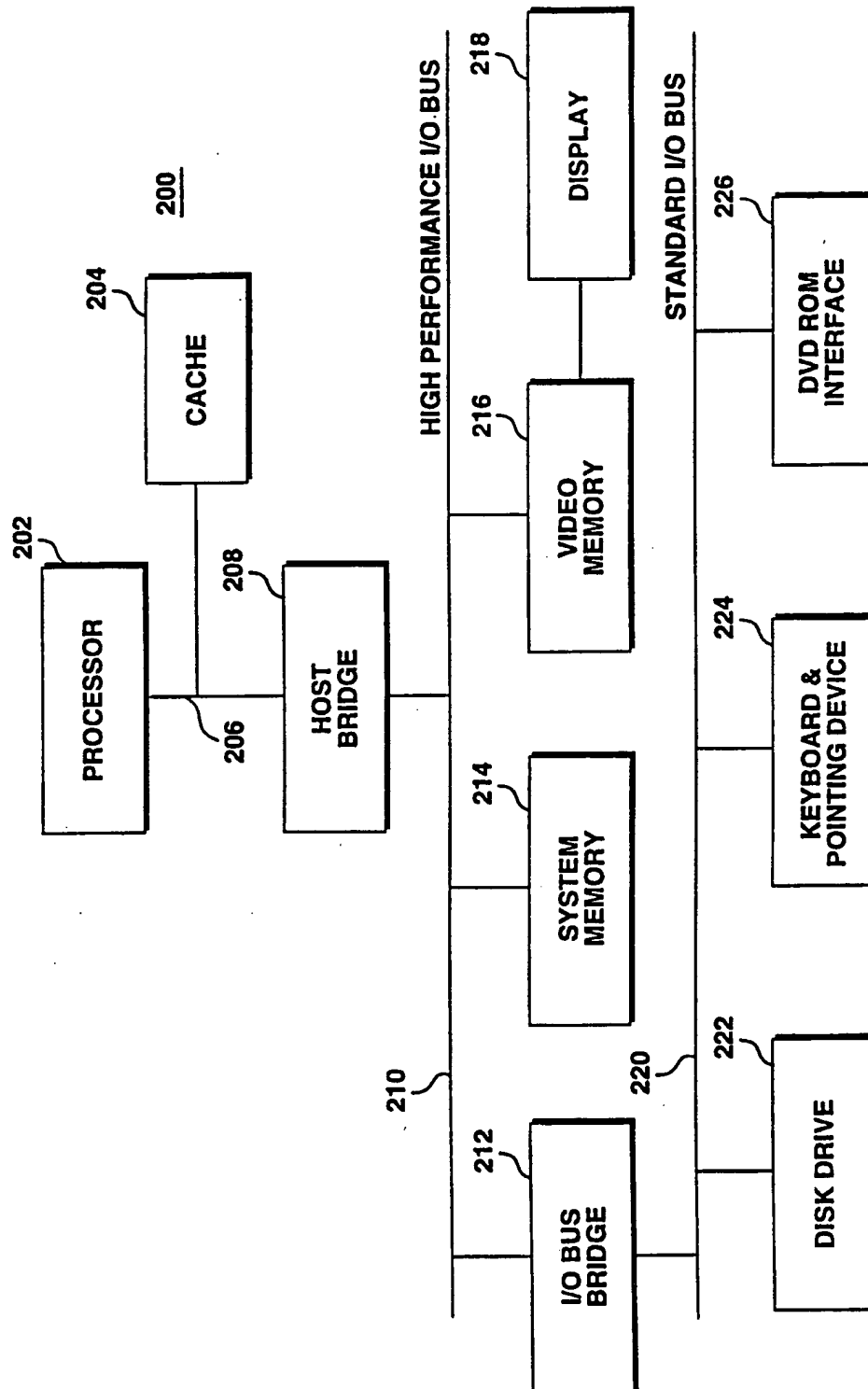


Fig. 9

INTERNATIONAL SEARCH REPORT

International application No.

PCT/US98/14914

A. CLASSIFICATION OF SUBJECT MATTER

IPC(6) : H04L 9/00

US CL : 395/186

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

U.S. : 395/186, 188.01; 380/3, 4, 9, 23, 25

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practicable, search terms used)

APS

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category*	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y	US 4,262,329 A (BRIGHT ET AL.) 14 APRIL 1981, see the abstract, col. 1, line 50 to col. 2, line 2.	1-31
Y	US 4,723,284 A (MUNCK ET AL.) 02 FEBRUARY 1988, see the abstract; also col. 2, lines 38-62.	1-31
Y	US 4,786,790 A (KRUSE ET AL.) 22 NOVEMBER 1988, see the abstract; also col. 1, line 52 to col. 3, line 48.	1-31
Y	US 5,224,160 A (PAULINI ET AL.) 29 JUNE 1993, see the abstract; col. 3, line 54 et seq.	1-311-



Further documents are listed in the continuation of Box C.



See patent family annex.

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A document defining the general state of the art which is not considered to be of particular relevance	*X* document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone
B earlier document published on or after the international filing date	*Y* document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art
L document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)	*A* document member of the same patent family
O document referring to an oral disclosure, use, exhibition or other means	
P document published prior to the international filing date but later than the priority date claimed	

Date of the actual completion of the international search

08 OCTOBER 1998

Date of mailing of the international search report

17 DEC 1998

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INTERNATIONAL SEARCH REPORT

International application No.
PCT/US98/14914

C (Continuation). DOCUMENTS CONSIDERED TO BE RELEVANT

Category*	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y	US 4,926,480 A (CHAUM) 15 MAY 1990, see the abstract, col. 4, line 50 to col. 5, line 34 and passim.	1-31
Y	US 5,265,164 A (MATYAS ET AL.) 23 NOVEMBER 1993, see the abstract.	1-31
Y	US 5,347,579 A (BLANDFORD) 13 SEPTEMBER 1994, see the abstract.	1-31
Y	US 5,081,675 A (KITIRUTSUNETORN) 14 JANUARY 1992, see the abstract; col. 2, line 1-12.	1-31
Y,P	US 5,768,382 A (SCHNEIER ET AL.) 16 JUNE 1998, see the abstract; figures 4A, 5.	1-31